Convex Quadratic Programming Problem

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Objective

• Understand the Quadratic Programming (QP)

Given a matrix $\mathbf{A} \in \mathbb{C}^{m \times n}$ of rank *r*, there exist unitary matrices $\mathbf{U} \in \mathbb{C}^{m \times m}$ and $V \in \mathbb{C}^{n \times n}$ such that

$$
\mathbf{A} = \mathbf{U} \Sigma \mathbf{V}^H, \text{ where } \Sigma = \begin{bmatrix} \mathbf{S} & 0 \\ 0 & 0 \end{bmatrix}_{m \times n}, \text{ and } \mathbf{S} = \text{diag}\{\sigma_1, \sigma_2, \dots, \sigma_r\},
$$

where $\sigma_1 > \sigma_2 > \cdots > \sigma_r > 0$.

- The matrix decomposition $\mathbf{A} = \mathbf{U} \Sigma \mathbf{V}^H$ is known as the **singular-value** decomposition (SVD) of **A**. $\{\cdot\}^H$ is a complex conjugate transpose. If **A** is a real-values matrix, then **U** and **V** become orthogoanl matrices and **V***^H* becomes **V***^T* .
- The positive scalars σ_i for $i = 1, 2, \ldots, r$ are called the **singular values** of **A**.
- $\mathbf{v} = \begin{bmatrix} u_1 & u_2 & \cdots & u_m \end{bmatrix}$ and $V = \begin{bmatrix} v_1 & v_2 & \cdots & v_n \end{bmatrix}$, vectors u_i and v_i are called the left and right singular vectors of **A**. We have

$$
\mathbf{A}\mathbf{A}^H = \mathbf{U} \begin{bmatrix} \mathbf{S}^2 & 0 \\ 0 & 0 \end{bmatrix}_{m \times m} \mathbf{U}^H, \quad \mathbf{A}^H \mathbf{A} = V \begin{bmatrix} \mathbf{S}^2 & 0 \\ 0 & 0 \end{bmatrix}_{n \times n} \mathbf{V}^H
$$

• See lecture11/svd_ex.jl. 3/29

The Moore-Penrose pseudo-inverse of a matrix $\mathbf{A} \in \mathbb{C}^{m \times n}$ is defined as the matrix $A^{\dagger} \in \mathbb{C}^{n \times m}$ that satisfies the following four conditions:

- 1. $\mathbf{A}\mathbf{A}^{\dagger}\mathbf{A} = \mathbf{A}$
- 2. $A^{\dagger} A A^{\dagger} = A^{\dagger}$
- 3. $(AA^{\dagger})^H = AA^{\dagger}$
- 4. $(\mathbf{A}^\dagger \mathbf{A})^H = \mathbf{A}^\dagger \mathbf{A}$

The example of **A***†* is (**A***^T* **A**)*−*1**A***^T* . The Moore-Penrose pseudo-inverse of **A** can be obtained as

$$
\mathbf{A}^{\dagger} = \mathbf{V} \Sigma^{\dagger} \mathbf{U}^{H}, \quad \Sigma^{\dagger} = \begin{bmatrix} \mathbf{S}^{-1} & 0 \\ 0 & 0 \end{bmatrix}_{n \times m}, \text{ and } \mathbf{S}^{-1} = \text{diag}\{\sigma_1^{-1}, \sigma_2^{-1}, \dots, \sigma_r^{-1}\}
$$

$$
\mathbf{A}^{\dagger} = \sum_{i=1}^{r} \frac{v_i u_i^H}{\sigma_i}
$$

The Moore-Penrose pseudo-inverse

For an underdetermined system of linear equations

$$
Ax = b \implies x = A^{\dagger}b + V_r \phi
$$

Given

$$
\mathbf{A} = \begin{bmatrix} 2.8284 & -1 & 1 \\ 2.8284 & 1 & -1 \end{bmatrix}, \quad \mathbf{A}^T \mathbf{A} = \begin{bmatrix} 16 & 0 & 0 \\ 0 & 2 & -2 \\ 0 & -2 & 2 \end{bmatrix} = \mathbf{V} \begin{bmatrix} 16 & 0 & 0 \\ 0 & 4 & 0 \\ 0 & 0 & 0 \end{bmatrix} \mathbf{V}^T
$$

The nonzero singular values of ${\bf A}$ are $\sigma_1 = \sqrt{16} = 4$ and $\sigma_2 = \sqrt{4} = 2$. By using a command $[UI, S1, V1] = svd(A'*A)$ of Matlab or U1, S1, V1 = svd $(A'*A)$ of Julia to obtain *S* and V= V1. Then using [U2, S2, V2] = svd(A) of Matlab or U2, S2, $V2 = svd(A)$ of Julia to get U=U2.

The Moore-Penrose pseudo-inverse

We obtain

$$
\mathbf{V} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0.7071 & -0.7071 \\ 0 & -0.7071 & -0.7071 \end{bmatrix} = \begin{bmatrix} v_1 & v_2 & v_3 \end{bmatrix},
$$

$$
\mathbf{U} = \begin{bmatrix} 0.7071 & -0.7071 \\ 0.7071 & 0.7071 \end{bmatrix} = \begin{bmatrix} u_1 & u_2 \end{bmatrix}
$$

$$
\mathbf{A}^{\dagger} = \frac{v_1 u_1^T}{\sigma_1} + \frac{v_2 u_2^T}{\sigma_2} = \begin{bmatrix} 0.1768 & 0.1768 \\ -0.2500 & 0.2500 \end{bmatrix}
$$

QP Problem

Consider a problem

minimize
$$
f(\mathbf{x}) = \frac{1}{2}\mathbf{x}^T H \mathbf{x} + \mathbf{x}^T \mathbf{p}
$$

subject to $\mathbf{A}\mathbf{x} = \mathbf{b}$, $\mathbf{A} \in \mathbb{R}^{p \times n}$

The solution of $\mathbf{A}\mathbf{x} = \mathbf{b}$ is $\mathbf{x} = \mathbf{V}_r\boldsymbol{\phi} + \mathbf{A}^\dagger\mathbf{b}$, where $\boldsymbol{\phi}$ is an arbitrary *r*-dimensional vector with $r = n - p$, \mathbf{A}^{\dagger} denotes the Moore-Penrose pseudo-inverse of \mathbf{A}, \mathbf{V}_r is composed of the last *r* columns of **V**, and $\mathbf{A} = \mathbf{U} \Sigma \mathbf{V}^T$. Then

$$
\underset{\phi}{\text{minimize}} = \frac{1}{2} \phi^T \hat{\mathbf{H}} \phi + \phi^T \hat{\mathbf{p}},
$$

where $\hat{\mathbf{h}} = \mathbf{V}_r^T \mathbf{H} \mathbf{V}_r$ and $\hat{\mathbf{p}} = \mathbf{V}_r^T (\mathbf{H} \mathbf{A}^\dagger \mathbf{b} + \mathbf{p})$ The global minimizer of the problem is $\mathbf{x}^* = \mathbf{V}_r\boldsymbol{\phi}^* + \mathbf{A}^\dagger\mathbf{b}$, where $\boldsymbol{\phi}^*$ is the solution of the linear system of equations $\hat{\mathbf{H}}\boldsymbol{\phi} = -\hat{\mathbf{p}}$.

QP Problem: Example

Solve the QP problem:

minimize
$$
f(\mathbf{x}) = \frac{1}{2} (\mathbf{x}_1^2 + \mathbf{x}_2^2) + 2\mathbf{x}_1 + \mathbf{x}_2 - \mathbf{x}_3
$$

subject to $\mathbf{A}\mathbf{x} = \mathbf{b}$

where $\mathbf{A} = \begin{bmatrix} 0 & 1 & 1 \end{bmatrix}$, and $\mathbf{b} = 1$.

$$
\mathbf{V}_r = \begin{bmatrix} -0.7071 & -0.707 \\ 0.5 & -0.5 \\ -0.5 & 0.5 \end{bmatrix}, \quad \mathbf{A}^T \left(\mathbf{A}^T \mathbf{A} \right)^{-1} = \begin{bmatrix} 0 & 0.5 & 0.5 \end{bmatrix}^T,
$$

$$
\hat{\mathbf{H}} = \mathbf{V}_r^T \mathbf{H} \mathbf{V}_r = \begin{bmatrix} 1.5 & 0.5 \\ 0.5 & 1.5 \end{bmatrix}, \quad \hat{\mathbf{p}} = \mathbf{V}^T \left(\mathbf{H} \mathbf{A}^\dagger \mathbf{b} + \mathbf{p} \right) = \begin{bmatrix} -0.2322 \\ -3.7677 \end{bmatrix}
$$

$$
\mathbf{x}^* = \mathbf{V}_r \boldsymbol{\phi}^* + \mathbf{A}^\dagger \mathbf{b} = \begin{bmatrix} -2 & -2 & 3 \end{bmatrix}^T
$$

The quadratic problem is

minimize
$$
\frac{1}{2} \mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{p}^T \mathbf{x} + \mathbf{c}
$$

subject to $\mathbf{A} \mathbf{x} = \mathbf{b}$
 $\mathbf{C} \mathbf{x} \le \mathbf{d}$

Example: least-squares

minimize $||Ax - b||_2^2$

- analytical solution **x** *∗* = **A***†***b**
- we can add linear constraints , e.g. *l ≤* **x** *≤ u*

Consider

minimize
$$
f(\mathbf{x}) = \frac{1}{2}\mathbf{x}^T \mathbf{H}\mathbf{x} + \mathbf{x}^T \mathbf{p}
$$

subject to $\mathbf{A}\mathbf{x} = \mathbf{b}$

Assume $\mathbf{A} \in \mathbb{R}^{p \times n}$ is of full row rank and $p < n$. By using the first-order necessary conditions $\nabla f_0(\mathbf{x}^*) + \mathbf{A}^T \boldsymbol{\lambda} = 0$, we have

$$
\mathbf{Hx}^* + \mathbf{p} + \mathbf{A}^T \mathbf{\lambda}^* = 0
$$

$$
\mathbf{Ax}^* - \mathbf{b} = 0
$$

$$
\begin{bmatrix} \mathbf{H} & \mathbf{A}^T \\ \mathbf{A} & 0 \end{bmatrix} \begin{bmatrix} \mathbf{x}^* \\ \mathbf{\lambda}^* \end{bmatrix} = \begin{bmatrix} -\mathbf{p} \\ \mathbf{b} \end{bmatrix} \quad (1)
$$

If **H** is positive definite and **A** is of full row rank, then the system matrix (1) is nonsingular and the solution **x** *∗* is the unique global minimizer of the problem.

The block matrix inversion lemma is

$$
\begin{bmatrix} \mathbf{A} & \mathbf{D} \\ \mathbf{C} & \mathbf{B} \end{bmatrix}^{-1} = \begin{bmatrix} \Delta^{-1} & -\Delta^{-1} \mathbf{D} \mathbf{B}^{-1} \\ -\mathbf{B}^{-1} \mathbf{C} \Delta^{-1} & \mathbf{B}^{-1} + \mathbf{B}^{-1} \mathbf{C} \Delta^{-1} \mathbf{D} \mathbf{B}^{-1} \end{bmatrix}
$$

where

$$
\Delta = \mathbf{A} - \mathbf{D} \mathbf{B}^{-1} \mathbf{C}
$$

Then we have

$$
\begin{aligned} \boldsymbol{\lambda}^* &= -(\mathbf{A}\mathbf{H}^{-1}\mathbf{A}^T)^{-1}(\mathbf{A}\mathbf{H}^{-1}\mathbf{p} + \mathbf{b}) \\ \mathbf{x}^* &= -\mathbf{H}^{-1}(\mathbf{A}^T\boldsymbol{\lambda}^* + \mathbf{p}) \end{aligned}
$$

Sequential Quadratic Optimization Problems (SQP)

Equality Constrained SQP

$$
\begin{aligned}\n\text{minimize} & \quad f(\mathbf{x}) = \mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{p}^T \mathbf{x} \\
\text{subject to} & \quad \mathbf{h}(\mathbf{x}) = 0\n\end{aligned}
$$

The Lagrangian is $\mathcal{L}(\mathbf{x},\boldsymbol{\lambda})=f(\mathbf{x})+\boldsymbol{\lambda}^T\mathbf{h}(x).$ At the minimum point, we have

$$
\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \lambda) = \nabla f(\mathbf{x}) + \mathbf{J}_h^T \lambda = 0
$$

$$
\nabla_{\lambda} \mathcal{L}(\mathbf{x}, \lambda) = h(\mathbf{x}) = 0
$$

The quadratic approximation:

$$
\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x} + \Delta \mathbf{x}, \lambda + \Delta \lambda) = \nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \lambda) + \nabla^2 \mathcal{L}(\mathbf{x}, \lambda) \Delta \mathbf{x} + \frac{\partial}{\partial \mathbf{x}} (\nabla_{\lambda} \mathcal{L}(\mathbf{x}, \lambda))^T \Delta \lambda
$$

$$
= \nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \lambda) + \mathbf{H}_{\mathcal{L}} \Delta \mathbf{x} + \mathbf{J}_{h}^T \Delta \mathbf{x}
$$

Sequential Quadratic Optimization Problems (SQP)

$$
\nabla_{\lambda} \mathcal{L}(\mathbf{x} + \Delta \mathbf{x}, \lambda + \Delta \lambda) = \nabla_{\lambda} \mathcal{L}(\mathbf{x}, \lambda) + \nabla_{\lambda} (\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \lambda)) \Delta \mathbf{x} + \nabla_{\mathbf{x}}^{2} \mathcal{L}(\mathbf{x}, \lambda) \Delta \lambda
$$

\n
$$
= h(\mathbf{x}) + \mathbf{J}_{h} \Delta \mathbf{x}
$$

\n
$$
\begin{bmatrix} \mathbf{H}_{\mathcal{L}} & \mathbf{J}_{h}^{T} \\ \mathbf{J}_{h} & 0 \end{bmatrix} \begin{bmatrix} \Delta \mathbf{x} \\ \Delta \lambda \end{bmatrix} = \begin{bmatrix} -\nabla_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \lambda) \\ -h(\mathbf{x}) \end{bmatrix}
$$

where

$$
\mathbf{x}_{k+1} = \mathbf{x}_k + \alpha_k \Delta \mathbf{x}_k
$$

$$
\boldsymbol{\lambda}_{k+1} = \boldsymbol{\lambda}_k + \Delta \boldsymbol{\lambda}_k
$$

By using slack variables,

minimize $f(\tilde{\mathbf{x}}) = \frac{1}{2}\tilde{\mathbf{x}}^T H \tilde{\mathbf{x}} + \tilde{\mathbf{x}}^T \mathbf{p}$ subject to $A\tilde{x} \leq b$

minimize
$$
f(\mathbf{x}) = \frac{1}{2}\mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{x}^T \mathbf{p}
$$

subject to $\mathbf{A}\mathbf{x} = \mathbf{b}$
 $-\mathbf{x} \leq 0$ (2)

We have $\mathcal{L}(\mathbf{x}, \lambda, \mu) = \frac{1}{2}\mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{x}^T \mathbf{p} - \mu^T \mathbf{x} + \lambda^T (\mathbf{A} \mathbf{x} - \mathbf{b})$ and $g(\lambda,\mu) = \inf_{\lambda,\mu} \mathcal{L}(x,\lambda,\mu)$ if and only if $\nabla \mathcal{L} = 0$ or $-A^T \lambda + \mu - Hx = p$ The dual problem becomes

maximize
$$
h(\mathbf{x}, \lambda, \mu) = -\frac{1}{2} \mathbf{x}^T \mathbf{H} \mathbf{x} - \lambda^T \mathbf{b}
$$

subject to $-\mathbf{A}^T \lambda + \mu - \mathbf{H} \mathbf{x} = \mathbf{p}$
 $\mu \ge 0$ (3)

The KKT conditions

$$
\mathbf{A}\mathbf{x} - \mathbf{b} = 0 \text{ for } \mathbf{x} \ge 0
$$

$$
-\mathbf{A}^T \mathbf{\lambda} + \boldsymbol{\mu} - \mathbf{H}\mathbf{x} - \mathbf{p} = 0 \text{ for } \boldsymbol{\mu} \ge 0
$$

$$
\mathbf{X}\boldsymbol{\mu} = 0,
$$
 (4)

where $X = diag\{x_1, x_2, ..., x_n\}$.

Let $W = \{x, \lambda, \mu\}$ be a feasible for the problems (2). The duality gap can be obtained for $\{{\bf x}, \boldsymbol{\lambda}, \boldsymbol{\mu}\}$ as

$$
\delta(\mathbf{x}, \lambda, \mu) = f(\mathbf{x}) - h(\mathbf{x}, \lambda, \mu) = \mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{x}^T \mathbf{p} + \lambda^T \mathbf{b}
$$

= $\mathbf{x}^T (-\mathbf{A}^T \lambda + \mu) + \lambda^T \mathbf{b} = \mathbf{x}^T \mu$ (5)

Setting $\mathbf{w}(\tau) = \{\mathbf{x}(\tau), \boldsymbol{\lambda}(\tau), \boldsymbol{\mu}(\tau)\}\$ that satisfies the KKT condition, the last line of (4) is changed to

$$
\mathbf{X}\boldsymbol{\mu} = \tau \mathbf{e}
$$

$$
\mathbf{e} = \begin{bmatrix} 1 & 1 & \cdots & 1 \end{bmatrix}^T \in \mathbb{R}^n
$$

$$
\delta[\mathbf{x}(\tau), \boldsymbol{\lambda}(\tau), \boldsymbol{\mu}(\tau)] = \mathbf{x}^T(\tau) \boldsymbol{\mu}(\tau) = n\tau
$$

Hence the duality gap approaches zero linearly as *τ →* 0.

Let $\mathbf{w}_k = {\mathbf{x}_k, \lambda_k, \mu_k}$ be such that \mathbf{x}_k is strictly feasible for the primal problem (2) and $\{\lambda_k, \mu_k\}$ is strictly feasible for the dual problem (3). The increment set is $\boldsymbol{\delta}_w = \{\boldsymbol{\delta}_x, \boldsymbol{\delta}_\lambda, \boldsymbol{\delta}_\mu\}$. We need $\mathbf{w}_{k+1} = \{\mathbf{x}_{k+1}, \boldsymbol{\lambda}_{k+1}, \boldsymbol{\mu}_{k+1}\} = \mathbf{w}_k + \boldsymbol{\delta}_w$ remains strictly feasible and approaches the central path.

If w_k satisfy the KKT and central path condition with τ_{k+1} , we have

$$
\mathbf{A}(\mathbf{x}_k + \boldsymbol{\delta}_x) - \mathbf{b} = 0
$$

$$
-\mathbf{A}^T(\boldsymbol{\lambda}_k + \boldsymbol{\delta}_\lambda) + (\boldsymbol{\mu}_k + \boldsymbol{\delta}_\mu) - \mathbf{H}(\mathbf{x}_k + \boldsymbol{\delta}_x) - \mathbf{p} = 0
$$

$$
(\mathbf{X} + \Delta \mathbf{X})(\boldsymbol{\mu}_k + \boldsymbol{\delta}_\mu) = \tau_{k+1} \mathbf{e}
$$

$$
\mathbf{A}\boldsymbol{\delta}_{x} = 0
$$

\n
$$
-\mathbf{H}\boldsymbol{\delta}_{x} - \mathbf{A}^{T}\boldsymbol{\delta}_{\lambda} + \boldsymbol{\delta}_{\mu} = 0
$$

\n
$$
\Delta \mathbf{X}\boldsymbol{\mu}_{k} + \mathbf{X}\boldsymbol{\delta}_{\mu} + \Delta \mathbf{X}\boldsymbol{\delta}_{\mu} = \tau_{k+1} - \mathbf{X}\boldsymbol{\mu}_{k}
$$

\n
$$
\mathbf{M}\boldsymbol{\delta}_{x} + \mathbf{X}\boldsymbol{\delta}_{\mu} = \tau_{k+1} - \mathbf{X}\boldsymbol{\mu}_{k}
$$
\n(6)

where $\Delta \mathbf{X} = \text{diag}\{(\boldsymbol{\delta}_x)_1, (\boldsymbol{\delta}_x)_2, \ldots, (\boldsymbol{\delta}_x)_n\}$, $\mathbf{M} = \text{diag}\{(\boldsymbol{\mu}_k)_1, (\boldsymbol{\mu}_k)_2, \ldots, (\boldsymbol{\mu}_k)_n\}$, and Δ **X** δ _{*µ*} is neglected.

Solving (6) we have

$$
\delta_{\lambda} = -Yy
$$

\n
$$
\delta_{x} = -\Gamma \mathbf{X} \mathbf{A}^{T} \delta_{\lambda} - y
$$

\n
$$
\delta_{\mu} = \mathbf{H} \delta_{x} + \mathbf{A}^{T} \delta_{\lambda}
$$
\n(7)

where $\Gamma = (\mathbf{M} + \mathbf{X}\mathbf{H})^{-1}$, $\mathbf{Y} = (\mathbf{A}\Gamma\mathbf{X}\mathbf{A}^T)^{-1}\mathbf{A}$ and $y = \Gamma(\mathbf{X}\boldsymbol{\mu}_k - \tau_{k+1}\mathbf{e})$

Since $\mathbf{x}_k > 0$ and $\boldsymbol{\mu}_k > 0$, matrices **X** and **M** are positive definite. Therefore **X***−*1**M** + **H** is also positive definite and the inverse of matrix

$$
\mathbf{M} + \mathbf{X} \mathbf{H} = \mathbf{X} (\mathbf{X}^{-1} \mathbf{M} + \mathbf{H})
$$

Since **A** is of full row rank, $\mathbf{A} \Gamma \mathbf{X} \mathbf{A}^T = \mathbf{A} (\mathbf{X}^{-1} \mathbf{M} + \mathbf{H})^{-1} \mathbf{A}^T$

Primal-dual path-following algorithm for the Convex QP problem

- 1. Input $\mathbf{A}, \mathbf{b}, \mathbf{c}$, and $\mathbf{w}_0 = \{\mathbf{x}_0, \boldsymbol{\lambda}_0, \boldsymbol{\mu}_0\}$. Set $k = 0$ and $\rho > \sqrt{n}$ (n is a dimension of **x**), and initialize the tolerance *ε* for the duality gap.
- 2. If $\boldsymbol{\mu}_k^T x_k ≤ ε$, output solution $\mathbf{w}^* = \mathbf{w}_k$ and stop; otherwise, continue with Step 3

3. Set
$$
\tau_{k+1} = \frac{\mu_k^T \mathbf{x}_k}{n+\rho}
$$
 and compute $\delta_w = \{\delta_x, \delta_\lambda, \delta_\mu\}$ using (7).

4. compute step size α_k as follow:

$$
\alpha_k = (1 - 10^{-6}) \alpha_{\text{max}} \quad \alpha_{\text{max}} = \min(\alpha_p, \alpha_d)
$$

where

$$
\alpha_p = \min_{i \text{ with } (\delta_x)_{i} < 0} \left[-\frac{(\mathbf{x}_k)_i}{(\delta_x)_i} \right], \quad \alpha_d = \min_{i \text{ with } (\delta_\mu)_i < 0} \left[-\frac{(\mu_k)_i}{(\delta_\mu)_i} \right]
$$

- The previous algorithm requires a strictly feasible **w**0, which might be difficult to obtain particularly for large-scale problems.
- Let $\mathbf{w}_k = {\mathbf{x}_k, \lambda_k, \mu_k}$ be such that $\mathbf{x}_k > 0$ and $\mu_k > 0$ but may not satisfy the central path condition. We need to find the next iterate

$$
\mathbf{w}_{k+1} = \mathbf{x}_k + \alpha_k \boldsymbol{\delta}_w
$$

such that $\mathbf{x}_{k+1} > 0$ and $\boldsymbol{\mu}_{k+1} > 0$ and that $\boldsymbol{\delta}_w = {\{\delta_x, \delta_\lambda, \delta_\mu\}}$ satisfies the equations

$$
-\mathbf{H}(x_k + \delta_x) - \mathbf{p} - \mathbf{A}^T(\lambda_k + \delta_\lambda) + (\mu_k + \delta_\mu) = 0
$$

$$
\mathbf{A}(\mathbf{x}_k + \delta_x) = \mathbf{b}
$$

$$
\mathbf{M}\delta_x + \mathbf{X}\delta_\mu = \tau_{k+1}\mathbf{e} - \mathbf{X}\mu_k
$$

$$
-\mathbf{H}\boldsymbol{\delta}_{x} - \mathbf{A}^{T}\boldsymbol{\delta}_{\lambda} + \boldsymbol{\delta}_{\mu} = \mathbf{r}_{d}
$$

$$
\mathbf{A}\boldsymbol{\delta}_{x} = \mathbf{r}_{p}
$$

$$
\mathbf{M}\boldsymbol{\delta}_{x} + \mathbf{X}\boldsymbol{\delta}_{\mu} = \tau_{k+1}\mathbf{e} - \mathbf{X}\boldsymbol{\mu}_{k}
$$

$$
\mathbf{r}_d = \mathbf{H}\mathbf{x}_k + \mathbf{p} + \mathbf{A}^T \mathbf{\lambda}_k - \boldsymbol{\mu}_k
$$

$$
\mathbf{r}_p = \mathbf{b} - \mathbf{A}\mathbf{x}_k
$$

The solution of **w** can be obtained as

$$
\delta_{\lambda} = -\mathbf{Y}_0 (\mathbf{A} \mathbf{y}_d + \mathbf{r}_p)
$$

\n
$$
\delta_x = -\Gamma \mathbf{X} \mathbf{A}^T \delta_{\lambda} - y_d \qquad (8)
$$

\n
$$
\delta_{\mu} = \mathbf{H} \delta_x + \mathbf{A}^T \delta_{\lambda} + \mathbf{r}_d
$$

$$
\Gamma = (\mathbf{M} + \mathbf{X}\mathbf{H})^{-1}
$$

$$
\mathbf{Y}_0 = (\mathbf{A}\Gamma\mathbf{X}\mathbf{A}^T)^{-1}
$$

$$
y_d = \Gamma[\mathbf{X}(\boldsymbol{\mu}_k + \mathbf{r}_d) - \tau_{k+1}\mathbf{e}]
$$

Nonfeasible-initialization Primal-dual path-following algorithm for the Convex QP problem

- 1. Input $\mathbf{A}, \mathbf{b}, \mathbf{c}$, and $\mathbf{w}_0 = \{\mathbf{x}_0, \boldsymbol{\lambda}_0, \boldsymbol{\mu}_0\}$. Set $k = 0$ and $\rho > \sqrt{n}$ (n is a dimension of *x*), and initialize the tolerance *ε* for the duality gap.
- 2. If $\boldsymbol{\mu}_k^T {\bf x}_k \leq \varepsilon$, output solution ${\bf w}^* = {\bf w}_k$ and stop; otherwise, continue with Step 3
- 3. Set $\tau_{k+1} = \frac{\mu_{k}^{T}x_{k}}{n+\rho}$ and compute $\pmb{\delta}_{w} = \{\pmb{\delta}_{x},\pmb{\delta}_{\lambda},\pmb{\delta}_{\mu}\}$ using (8)
- 4. compute step size α_k as follow:

$$
\alpha_k = (1 - 10^{-6})\alpha_{\text{max}} \quad \alpha_{\text{max}} = \min(\alpha_p, \alpha_d)
$$

where

$$
\alpha_p = \min_{i \text{ with } (\delta_x)_i < 0} \left[-\frac{(\mathbf{x}_k)_i}{(\delta_x)_i} \right], \quad \alpha_d = \min_{i \text{ with } (\delta_\mu)_i < 0} \left[-\frac{(\mu_k)_i}{(\delta_\mu)_i} \right]
$$
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Solve the convex QP problem

minimize
$$
f(\mathbf{x}) = \frac{1}{2}x^T \begin{bmatrix} 4 & 0 & 0 \\ 0 & 1 & -1 \\ 0 & -1 & 1 \end{bmatrix} \mathbf{x} + \mathbf{x}^T \begin{bmatrix} -8 \\ -6 \\ -6 \end{bmatrix}
$$

subject to $x_1 + x_2 + x_3 = 3$
 $\mathbf{x} \ge 0$

- Using qp path sf.m with $\mathbf{x}_0 = [1 \ 1 \ 1]^T$, $\boldsymbol{\lambda}_0 = 7$, $\boldsymbol{\mu}_0 = [3 \ 1 \ 1]^T$, which is combined as a strictly feasible point w_0 . Using 14 iterations, the solution is $\mathbf{x}^* = [0.50 \ 1.25 \ 1.25]^T$.
- Using gp_path_nf.m with $\mathbf{x}_0 = [1 \ 2 \ 2]^T$, $\boldsymbol{\lambda}_0 = 1$, $\boldsymbol{\mu}_0 = [0.2 \ 0.2 \ 0.2]^T$, which is not a strictly feasible point w_0 . Using 16 iteration, the solution is $\mathbf{x}^* = [0.50 \ 1.25 \ 1.25]^T$.

Solve the shortest-distance between triangle *R* and *S* shown in Figure below and the point *r ∗ ∈ R* and *s ∗ ∈ S* that yield the minimum distance. 1 2 3 *r ∗ s ∗ R S x*1*, x*³ *x*2*, x*⁴

Solution: Let $\mathbf{r} = [x_1 \ x_2]^T \in \mathcal{R}$ and $\mathbf{s} = [x_3 \ x_4]^T \in \mathcal{S}$. The square of the distance between *r* and *s* is given by

$$
(x_1 - x_3)^2 + (x_2 - x_4)^2 = \mathbf{x}^T \mathbf{H} \mathbf{x}, \qquad \mathbf{H} = \begin{bmatrix} 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \\ -1 & 0 & 1 & 0 \\ 0 & -1 & 0 & 1 \end{bmatrix}, \quad \mathbf{x} = \begin{bmatrix} x_1 & x_2 & x_3 & x_4 \end{bmatrix}^T
$$

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 Ω

1 2 3

The constraints of this problem are

*x*₁ ≥ 0*, x*₂ ≥ 0*, x*₁ + 2*x*₂ ≤ 2*, x*₄ ≥ 2*, x*₃ + *x*₄ ≥ 3*, x*₃ + 2*x*₄ ≤ 6

The problem can be formulated as the QP problem

minimize
$$
f(\mathbf{x}) = \frac{1}{2}\mathbf{x}^T \mathbf{H} \mathbf{x}
$$

\nsubject to $\mathbf{A}\mathbf{x} \leq b$
\n
$$
\mathbf{A} = \begin{bmatrix}\n-1 & 0 & 0 & 0 \\
0 & -1 & 0 & 0 \\
1 & 2 & 0 & 0 \\
0 & 0 & 0 & -1 \\
0 & 0 & 1 & 2\n\end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix}\n0 \\
0 \\
2 \\
-2 \\
-3 \\
6\n\end{bmatrix}
$$

• We need to convert the problem into the form of (3). By letting **x** = **x** ⁺ *−* **x***[−]* where $\mathbf{x}^+\geq 0$ and $\mathbf{x}^-\geq 0$, and then introducing slack vector $\eta\geq 0.$

Changing to Standard form, we have

$$
\frac{1}{2}\mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{x}^T \mathbf{p} \qquad \Rightarrow \frac{1}{2}(\mathbf{x}^+ - \mathbf{x}^-)^T \mathbf{H} (\mathbf{x}^+ - \mathbf{x}^-) + (\mathbf{x}^+ - \mathbf{x}^-)^T \mathbf{p}
$$

$$
\mathbf{A} \mathbf{x} \le \mathbf{b} \qquad \Rightarrow \mathbf{A} (\mathbf{x}^+ - \mathbf{x}^-) + \mathbf{\eta} = \mathbf{b}
$$

Then the problem is changed to

minimize
$$
\frac{1}{2}\hat{\mathbf{x}}^T \hat{\mathbf{H}} \hat{\mathbf{x}} + \hat{\mathbf{x}}^T \hat{\mathbf{p}}
$$

\nsubject to $\hat{\mathbf{A}} \hat{\mathbf{x}} = \mathbf{b}$
\n $\hat{\mathbf{x}} \ge 0$
\n $\hat{\mathbf{H}} = \begin{bmatrix} \mathbf{H} & -\mathbf{H} & 0_{4 \times 6} \\ -\mathbf{H} & \mathbf{H} & 0_{4 \times 6} \\ 0_{6 \times 4} & 0_{6 \times 4} & 0_{6 \times 6} \end{bmatrix}, \quad \hat{\mathbf{p}} = \begin{bmatrix} \mathbf{p} \\ -\mathbf{p} \\ 0_{6 \times 1} \end{bmatrix}, \quad \hat{\mathbf{x}} = \begin{bmatrix} \mathbf{x}^+ \\ \mathbf{x}^- \\ \eta \end{bmatrix}, \quad \hat{\mathbf{A}} = \begin{bmatrix} \mathbf{A} & -\mathbf{A} & \mathbf{I}_{6 \times 6} \end{bmatrix}$

- We use nonfeasible-initialization method with $\mathbf{x}_0 = \text{ones}\{14, 1\}$, λ_0 = ones $\{6, 1\}$, μ_0 = ones $\{14, 1\}$, where ones $\{m, 1\}$ represents a colmun vector of dimension *m* whose elements are all equal to one.
- \cdot Setting $\varepsilon = 10^{-5}$ and $\rho = n + 20\sqrt{n}$, the algorithm using gp _path_nf.m takes 11 iterations to converge to *x*ˆ *∗*.
- The solution of **x** is

$$
\mathbf{x}^* = \hat{\mathbf{x}}^*[1:4] - \hat{\mathbf{x}}^*[5:8] = \begin{bmatrix} 0.4\\0.8\\1.0\\2.0 \end{bmatrix}
$$

• The shortest distance is

shortest distance =
$$
\sqrt{(\mathbf{x}^*)^T \mathbf{H} \mathbf{x}^*}
$$
 = 1.3416 unit

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